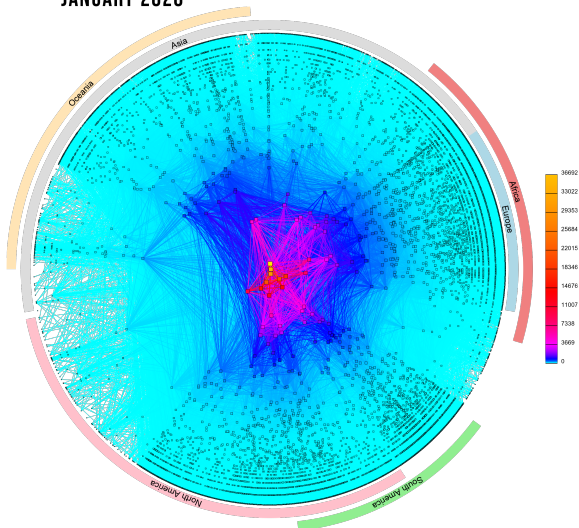


UE PRIP  
Principes des réseaux informatiques par la  
pratique  
**Introduction**

Isabel Amigo

2022

# CAIDA'S IPV4 AS CORE GRAPH JANUARY 2020



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source <https://www.caida.org/projects/as-core/2020/>

# Today's objective

- get “feel” and some terminology
- more depth and detail later in course

## Approach:

1. Introductory course using Internet as example
2. Individual quizz
3. Labs/TPs

# Today's roadmap

1. what's the Internet?
2. what's a protocol?
3. network edge: hosts, access net, physical media
4. network core: packet/circuit switching
5. protocol layers, service models, encapsulation
6. Internet's structure

# 1. What's the Internet?

# What's the Internet?

- A. “Construction” point of view (describing elements)
- B. Service point of view

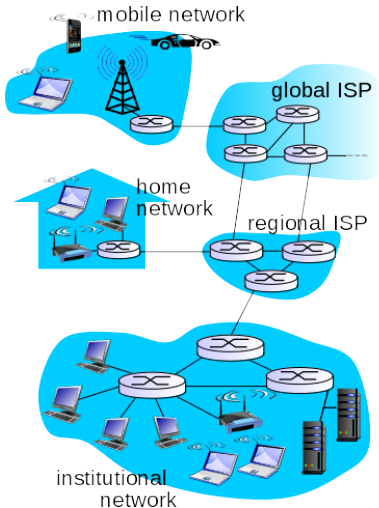
# What's the Internet?

A. "Construction" point of view (describing elements)

- millions of connected computing devices: hosts running network apps
- communication links: fiber, copper, radio, satellite  
transmission rate: *bandwidth*
- routers: forward packets (chunks of data)

# What's the Internet?

A. "Construction" point of view (describing elements)





# What's the Internet?

A. A “construction” point of view (describing elements)

Some characteristics:

- Internet: “network of networks” loosely hierarchical
- protocols e.g., TCP, IP, HTTP, Skype, Ethernet
- Internet standards
  - RFC: Request for comments
  - IETF: Internet Engineering Task Force
  - IEEE

Some questions:

**Q** public or private?

**Q** who governs Internet?

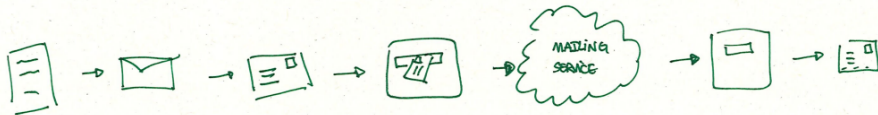
# What's the Internet?

## B. A services view

- A “thing” that allows us to use distributed applications
  - Q where do these applications run?
  - Q what are these applications?
- To use this “thing”, we have to respect some “rules ”
  - analogy: mail (physical)
  - APIs and Protocols (next slides)

thing = communication infrastructure

“rules” = APIs and protocols



## 2. What's a protocol?

# What's a protocol?

Human protocols:

- introductions “what’s the time?” “I have a question”
- specific msgs sent
- specific actions taken when msgs received, or other events

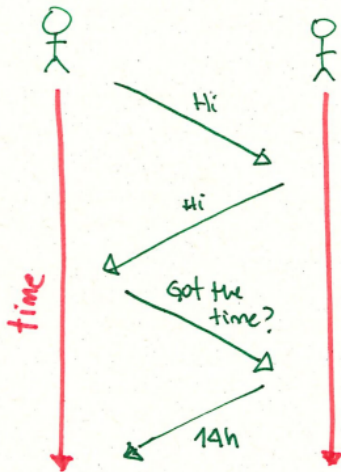
Network protocols:

- analogue to human protocols but
- machines rather than humans
- all communication activity in Internet governed by protocols

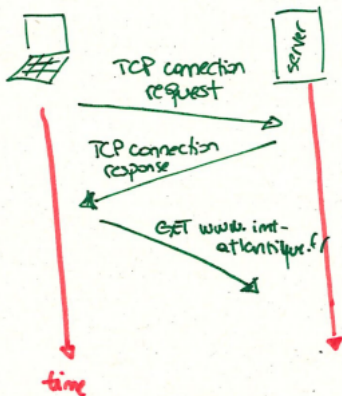
# What's a protocol?

## Examples

### HUMAN protocol



### COMPUTER protocol

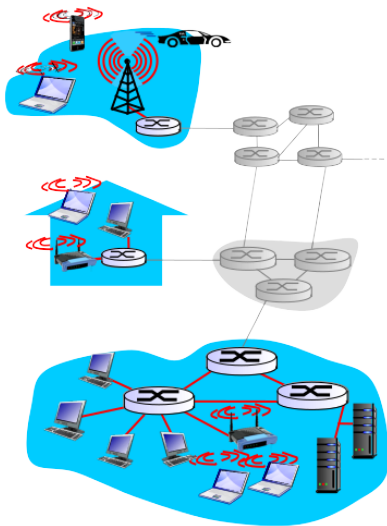


# What's a protocol?

Protocols define format, order of msgs sent and received among network entities, and actions taken on msg transmission and reception.

### 3. Network edge

# Network Edge



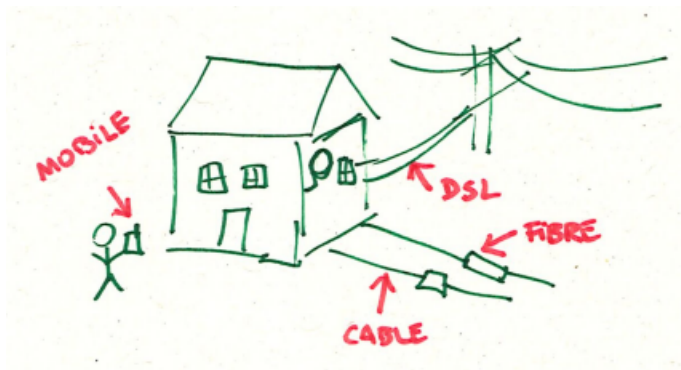


# Network Edge

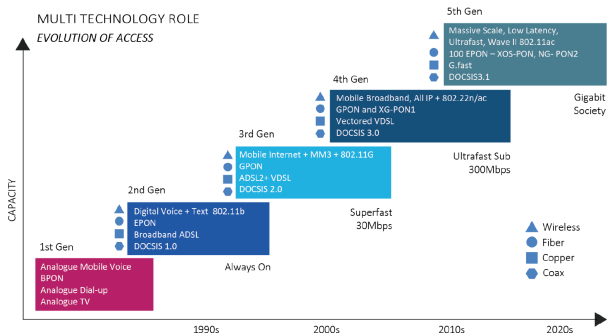
What is in the networks' edge?

- end systems
  - run applications (e.g. email, web,...)
- access networks
  - residential
  - institutional
  - mobile
- access technologies
  - physical media: wireless, wired, fiber
  - evolutions e.g. 1G, 2G .. VDSL, VDSL2, IEEE 802.11 a, b ...  
ax

# Broadband access technologies



# Broadband access technologies



Source: ADTRAN relayed by <http://reports.broadbandcommission.org/>

Keep in mind! Bandwidth (bps)? dedicated or shared?

# Physical media

Copper, Coax, Fiber, Air



Twisted Pair (copper), RJ45 connector for twisted pair, RJ45 and Coaxial connectors, Ethernet NIC (RJ45), Coaxial NIC, Fiber connectors

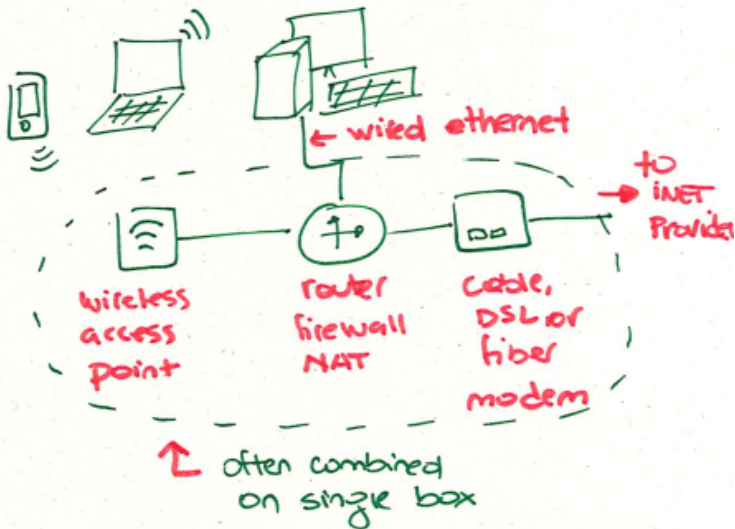
# Access networks

## Mobile access



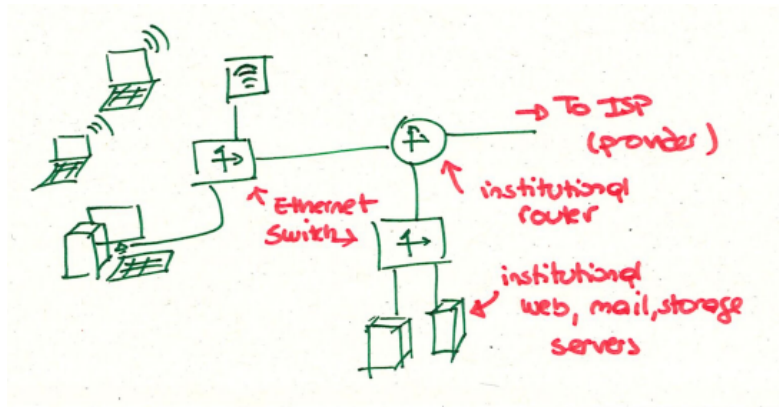
# Access networks

Home/ Residential network



# Access networks

Institutional network



# End hosts

End host:

- takes applications' message and breaks it into chunks of data (**packets**) of say length  $L$  bits
- sends the packet through a network interface at **transmission rate**  $R$  bps
  - transmission rate also known as bandwidth also known as capacity

**Transmission** delay ( $D$ ) is thus given by  $L$  and  $R$ .

$$D = \frac{L \text{ (bits)}}{R \text{ (bits/second)}}$$

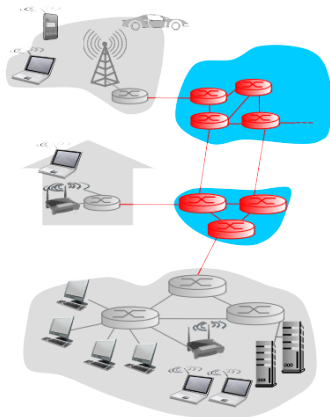
! There are other sources of end-to-end delay: propagation, queuing, processing, we'll see it later!



## 4. Network Core

# Network Core

Packet switching core



- we usually talk of aggregation network and core network, but today we are **simplifying**
- mesh of interconnected *routers*

## 4. Network Core

Routers forward packets from *one router to the next*, across links based on packet destination, each packet transmitted at full link capacity

# How is the *next router* determined?

Some key vocabulary: routing and forwarding

**Routing** function: determines source-destination route taken by packets on a *per-hop* basis

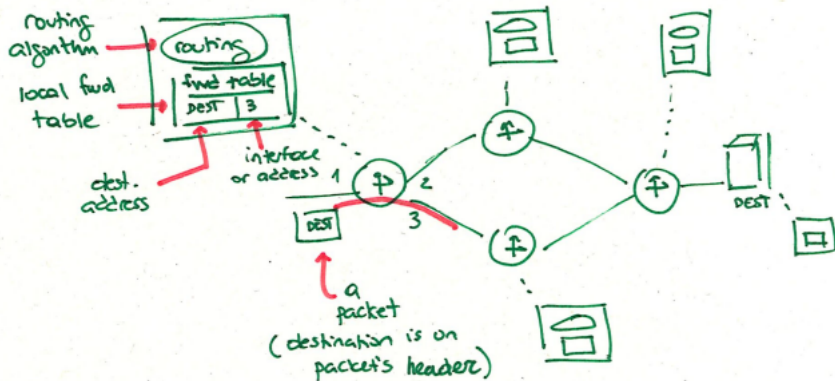
- The routing function, thanks to a routing algorithm, determines the routing table.

**Forwarding** function: moves packets from router's input interface to appropriate router outputs interface

- The forwarding function, uses such table to decide through which output interface send the packet

**!** Both functions performed at each router.

# Routing and Forwarding functions



# Store-and-forward packet-switching

Store and forward: entire packet must arrive at router before it can be transmitted on next link

- ! introduces a delay of  $D = L/R$  seconds at each *hop* (where as before  $L$  is the length of the packets in bits and  $R$  the interface's capacity in bps)
- ! This is how Internet works nowadays

# Alternative Core: Virtual circuit

End-end resources allocated to, reserved for, “call” between source, dest.

- dedicated resources: no sharing
  - circuit-like (guaranteed) performance
  - resource piece remains idle if not used by owning call (no sharing)
- dividing link bandwidth into “pieces”
  - frequency division
  - time division

Example: telephony network

# Circuit-switched vs Packet switching networks

Analogy: restaurant with and without reservations.

Restaurant with reservations:

- 😊 You arrive and are served immediately
- 😞 You have to call and reserve before going
- 😞 If you don't show, the restaurant wastes a table

Restaurant without reservations:

- 😊 You might have to wait on arrival
- 😊 You don't have to call in advance to reserve
- 😊 Restaurant doesn't take the risk of wasting a table



# Statistical multiplexing

Packet switching allows more users to use network!

## Example

- 1 Mbps link
- each user is active (i.e. sends data) 10% of the time
- each user when active sends data at rate 100kbps

**Q** How many users can be supported?

- Circuit-switching: 10 users
- Packet-switching: With 35 users probability of more than 10 users transmitting at the same time is  $\sim 0.0004$

**Q** How did we get value 0.0004?

**Q** What happens if there are more than 35 users?

# Is packet-switching *better* than circuit-switching?

☺ Packet-switching is great for bursty data

- resource sharing
- simpler, no call setup

☹ No performance guarantees, excessive congestion possible:  
packet delay and loss

- protocols needed for reliable data transfer and congestion control

**Q** Examples of applications generating bursty and non-bursty data?

**Q** How to provide bandwidth guarantees, needed for some applications, on packet-switching networks?

- several “patches” exist, but still a research problem!

# A word on end-to-end packet delay

Packet-switched networks experience different source of delays:

- **Transmission** at each node
- **Propagation** through physical media
- **Queuing** at nodes
- **Processing** at nodes

## 5. Layers model

# Protocol Layers

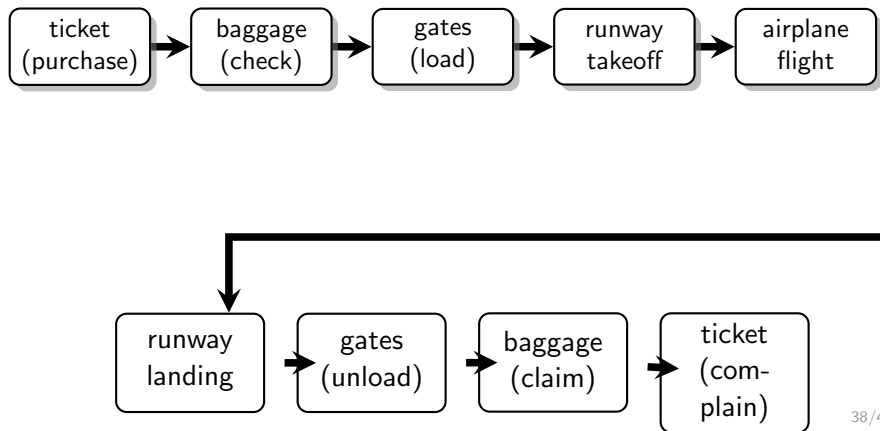
Networks are complex!

**Q** How can we organize the discussion to make it simpler?

**Q** How can we organize such systems to make them simpler, easier to develop, as little as possible prone to errors, upgradeable, scalable, ...?

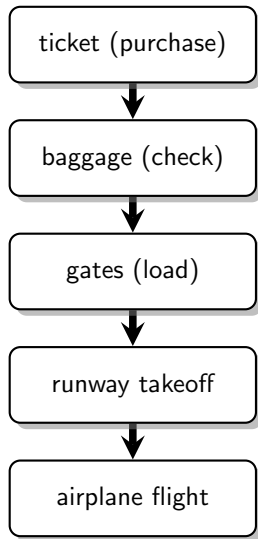
# Analogy: organization of a travel

Different steps



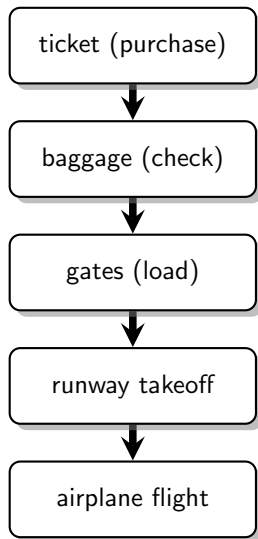
# Analogy: organization of a travel

Let's arrange steps differently



# Analogy: organization of a travel

Process on different entities but of same 'layer' can 'understand' each other





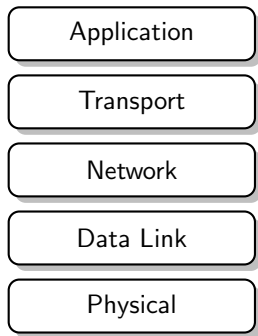
# Analogy: organization of a travel

We have constructed a layered service architecture

Each layer provides a service

- via its own internal-layer actions and functions
- relying on services provided by layer below
- each layer 'talks' to its homologue

# Layered reference model in the Internet



Layer 5

supporting network applications e.g. FTP, SMTP, HTTP

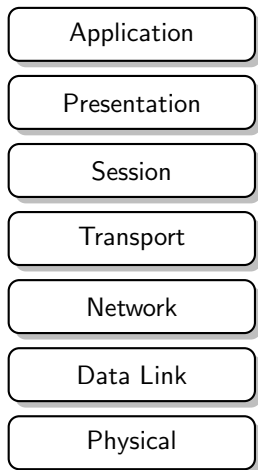
Layer 4

process-process data transfer e.g. TCP, UDP

Layer 3

routing of data across networks to destination IP address

# ISO/OSI reference model



! Internet stack “missing” two layers!  $\Rightarrow$  these services, if needed, must be implemented in application

**Q** Are they needed?

# Layers as a means to deal with complex systems

Some motivations for layering:

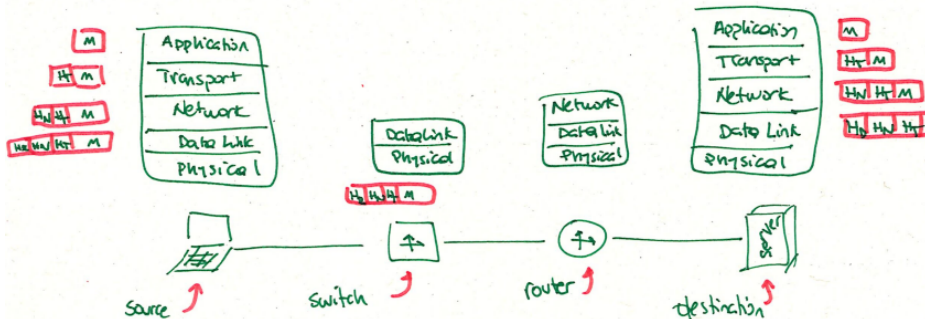
- ease implementation: focus on one function, rely on others for solving other 'problems'
- layered reference model for discussion
- modularization eases maintenance, updating of system
- change of implementation of layer's service transparent to rest of system e.g., change in gate procedure doesn't affect rest of system

**Q** Can you think of any cons of layering?

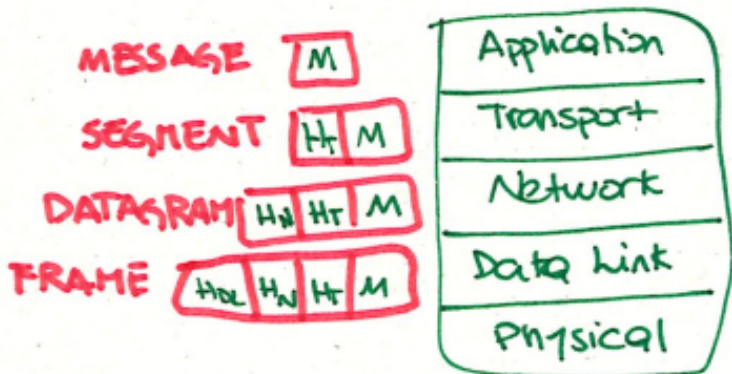
# Encapsulation

How does each process “talk” to its homologue on another equipment?

Messages are *encapsulated* i.e. extra information (known as *headers*) are added by each layer in the source, and interpreted by corresponding layers in middle equipments (e.g. switches, routers) and by the destination.



# Some vocabulary



**Q** Analogy between headers, encapsulation and our travel planning example?

# Summary

Make sure you have some initial idea of the following concepts!  
(we'll come back to them during next classes)

- Access network, Access technologies, Core network
- Packet
- Protocol
- Layers
- Delay

# Acknowledgements

The contents of these slides are partially taken from Computer Networking a Top Down approach, J. KUROSE and K. ROSS and K. KUROSE's networking course

[http://www-net.cs.umass.edu/cs453\\_fall\\_2013/](http://www-net.cs.umass.edu/cs453_fall_2013/)